Technische Universität München — *Dr. K. N. Verma*
Fakultät für Informatik — *verma@in.tum.de*
Prof. Dr. H. Seidl — *Room: MI 02.07.041*

Abstract Machines

Summer Semester 2004

5. Homework — *Deadline: 2 June 2004 12:00*

Exercise 1: *10 Points*

Have a look at the code generated for the expression $e \equiv (a + a)$ with $\rho = \{a \mapsto (L, 1)\}$ and $kp = 1$. It was created using the Call by Need strategy.

$$\begin{array}{lclcll}
\mathrm{code}_V\ e\ \rho\ 1 & = & \mathrm{getvar\ a}\ \rho\ 1 & = & 1 & \mathrm{pushloc}\ 0 \\
 & & \mathrm{eval} & & 2 & \mathrm{eval} \\
 & & \mathrm{getbasic} & & 2 & \mathrm{getbasic} \\
 & & \mathrm{getvar\ a}\ \rho\ 2 & & 2 & \mathrm{pushloc}\ 1 \\
 & & \mathrm{eval} & & 3 & \mathrm{eval} \\
 & & \mathrm{getbasic} & & 3 & \mathrm{getbasic} \\
 & & \mathrm{add} & & 3 & \mathrm{add} \\
 & & \mathrm{mkbasic} & & 2 & \mathrm{mkbasic}
\end{array}$$

The eval instructions check whether the value of a has been computed. If not a still has to be evaluated. The second occurrence of eval in the above code is redundant, because the value of a is already known at this point.

The code generation functions can be modified such that redundant eval instructions are not generated any more. To do so, extend the code generation function for an expression e with an additional argument A. A collects the set of visible variables that are bound outside e and that have always been evaluated when reaching the code to be generated for e.

Thus the code generation scheme for variable access shall look as follows:

$$\mathrm{code}_V\ x\ \rho\ kp\ A = \begin{cases} \mathrm{getvar}\ x\ \rho\ kp & ,\ \text{if } x \in A \\ \\ \mathrm{getvar}\ x\ \rho\ kp & \\ \mathrm{eval} & ,\ \text{otherwise} \end{cases}$$

For example:

$$\begin{array}{lcl}
\mathrm{code}_V\ (e_1\ \Box_2\ e_2)\ \rho\ \mathrm{kp}\ A & = & \mathrm{code}_B\ e_1\ \rho\ \mathrm{kp}\ A \\
 & & \mathrm{code}_B\ e_2\ \rho\ (\mathrm{kp}+1)\ A \cup A[e_1] \\
 & & \mathrm{op}_2; \mathrm{mkbasic}
\end{array}$$

where $A[e_1]$ is the set of free variables in the expression e_1 which already must have been evaluated in order to evaluate e_1.

a) Define formally $A[e]$, where e is a PuF expression.

b) Modify the code generation functions for PuF expressions in order to get rid of redundant eval instructions.

Exercise 2: *3+7 Points*

Generate code for the following expressions:

a)
```
1 + let
        x = g + 10;
        y = x * 4
      in y * g
```
with $\rho = \{g \mapsto (G, 1)\}$ and $kp = 1$.

b)
```
letrec
fib = fn x => if x <= 1 then 1 else (fib (x-1)) + (fib (x-2))
in fib 4
```
with $\rho = \emptyset$ and $kp = 0$.