

Lösungsvorschläge der Wiederholungsklausur zu Einführung in die Informatik II

Aufgabe 1      **Java-GUI**

```
import java.awt.*;
import java.awt.event.*;

public class Main extends Frame implements ActionListener
{
    private Button button = new Button("Quadriere");
    private TextField value = new TextField();

    public Main()
    {
        add(value, BorderLayout.NORTH);
        add(button);
        button.addActionListener(this);
    }

    public void actionPerformed(ActionEvent e)
    {
        try
        {
            int x = Integer.parseInt(value.getText());
            value.setText((x * x) + "");
        }
        catch(Exception ex) {}
    }

    public static void main(String[] args)
    {
        Frame frm = new Main();
        frm.setVisible(true);
    }
}
```

**Aufgabe 2**      **Datalog**

- a)  $\text{sohn}(X,Y) \text{ :- maennlich}(X), \text{kind}(X,Y).$
- b)  $\text{mutter}(X,Y) \text{ :- weiblich}(X), \text{kind}(Y,X).$
- c)  $\text{nachfahre}(X,Y) \text{ :- kind}(X,Y).$   
 $\text{nachfahre}(X,Z) \text{ :- kind}(X,Y), \text{nachfahre}(Y,Z).$
- d)  $\text{gemeinsames\_kind}(X,Y) \text{ :- kind}(K,X), \text{kind}(K,Y).$
- e)

$$\frac{\text{kind}(\text{hieronymus}, \text{rosalinde}) \quad \frac{\text{kind}(\text{rosalinde}, \text{konstanze})}{\text{nachfahre}(\text{rosalinde}, \text{konstanze})}}{\text{nachfahre}(\text{hieronymus}, \text{konstanze})}$$

### Aufgabe 3      **Ocaml**

```
let rec map2 l1 l2 f =
  match (l1,l2) with
  (h1::r1,h2::r2) -> (f h1 h2)::(map2 r1 r2 f)
  | _ -> []

let zip l1 l2 = map2 l1 l2 (fun a b -> (a,b))

let v = zip [1;2;3] [4;5;6]

type node = int
type graph = node list array

let inverse g =
  let gi = Array.make (Array.length g) []
  in Array.iteri
    (fun i neighbours ->
      List.iter
        (fun neighbour -> gi.(neighbour) <- i::gi.(neighbour))
        neighbours)
    g; gi

let g = Array.of_list [[4];[0];[0;3];[1];[1;2;3]]
let g1 = inverse g
```

## Aufgabe 4      **Bipartiter graph**

```
type node = int
type graph = node list array

exception NotBipartite

let rec isIn i l = match l with [] -> false | h::r -> if i=h then true else isIn i r

let makeBipartition g =
  let visited = Array.make (Array.length g) false
  in
  let rec doit i forPart1 (part1,part2) =
    (* forPart is true if the current partition is partition no 1 *)
    let visitNeighbours =
      (* auxiliary function to start the traversal for the neighbour nodes *)
      List.fold_left (fun (part1,part2) j -> doit j (not forPart1) (part1,part2))
    in
    if i>= Array.length g then (part1,part2)
    else
      if forPart1 then
        (* the current node is supposed to be in the first partition *)
        if (isIn i part2) then raise (NotBipartite)
      else
        if visited.(i) then (* the node was already visited *) (part1,part2)
        else (visited.(i) <- true; visitNeighbours (i::part1,part2) g.(i))
      else
        (* the current node is supposed to be in the second partition *)
        if (isIn i part1) then raise (NotBipartite)
    else if visited.(i) then (part1,part2)
    else (visited.(i) <- true; visitNeighbours (part1,i::part2) g.(i))
  in
  let rec visitAll i (part1,part2) =
    (* auxiliary function necessary if the graph is not connected *)
    if i>= Array.length g then (part1,part2)
    else if visited.(i) then visitAll (i+1) (part1,part2)
    else visitAll (i+1) (doit i true (part1,part2))
  in visitAll 0 ([],[])

let g = Array.of_list [[1;4];[2];[5];[];[3];[3];[7];[6]]
let g1 = makeBipartition g

(* Alternative *)

let bipart g =
  let nodeCount = Array.length g in
  let visited = Array.make nodeCount (-1) in
  let partitions = Array.make 2 [] in
  let rec dfs p n =
    if visited.(n) = -1 then
      begin
        visited.(n) <- p;
        partitions.(p) <- n::partitions.(p);
        List.iter (dfs ((p+1) mod 2)) g.(n)
      end
  in
```

```
end
else if visited.(n) != p then
  raise NotBipartite
in
for n = 0 to nodeCount - 1 do
  if visited.(n) = -1 then
    dfs 0 n
done;
(partitions.(0),partitions.(1))
```

### Aufgabe 5 Verifikation eines Min-Java-Programms (Lösungsvorschlag)

Für die Schleifen-Invariante raten wir:

$$F \equiv y = (1+x)^i \wedge xz = y-1$$

Dann ergibt sich:

$$D \equiv WP[i == n](F, E) \equiv (i \neq n \wedge F) \vee (i = n \wedge E) \equiv (i \neq n \wedge F) \vee (i = n \wedge F) \equiv F$$

$$H \equiv WP[i = i + 1](D) \equiv y = (1+x)^{i+1} \wedge xz = y-1$$

$$G \equiv WP[y = y * (x + 1)](H) \equiv y(x+1) = (1+x)^{i+1} \wedge xz = y(x+1) - 1$$

$$C \equiv WP[i = 0](D) \equiv y = 1 \wedge xz = y-1$$

$$B \equiv WP[y = 1](C) \equiv 1 = 1 \wedge xz = 1 - 1 = 0 \equiv xz = 0$$

$$A \equiv WP[z = 0](B) \equiv 0 = 0 \equiv true$$

Schlussendlich bleibt noch festzustellen, dass gilt:

$$WP[z = z + y](G) \equiv y(x+1) = (1+x)^{i+1} \wedge x(z+y) = y(x+1) - 1$$

$$\equiv y(x+1) = (1+x)^{i+1} \wedge xz = y-1$$

$$\leftarrow y = (1+x)^i \wedge xz = y-1 \equiv F$$

## Aufgabe 6      **Verifikation funktionaler Programme**

Zu zeigen ist das Prädikat

$$\text{flip}(\text{flip list}) = \text{list}$$

Der Beweis erfolgt durch Induktion über die Länge der Liste *list*. Die Länge der Liste *list* bezeichnen wir im Folgenden mit *n*.

**Induktionsanfang** ( $n = 0$ ):

$$\begin{aligned} &= \text{flip } [] \\ \stackrel{\text{Def.}}{=} & (\text{fun list} \rightarrow \text{match list with} \\ & \quad [] \rightarrow [] \\ & \quad | (x,y)::\text{tail} \rightarrow (y,x)::(\text{flip tail})) [] \\ &= \text{match } [] \text{ with} \\ & \quad [] \rightarrow [] \\ & \quad | (x,y)::\text{tail} \rightarrow (y,x)::(\text{flip tail}) \\ &= [] \end{aligned}$$

Da aus  $n = 0$  folgt, dass  $\text{list} = []$  ergibt sich:

$$\text{flip}(\text{flip list}) = \text{flip}(\text{flip } []) = \text{flip } [] = [] = \text{list}$$

**Induktionsschluß:** ( $n \rightarrow n + 1$ ) : Die Liste *list* hat die Länge  $n + 1$ . Daraus folgt, dass eine Liste *list'* der Länge  $n$  und ein Paar  $(x', y')$  existiert mit  $\text{list} = (x', y')::\text{list}'$ .

Weiterhin ergibt sich:

$$\begin{aligned} &= \text{flip } (a,b)::l \\ \stackrel{\text{Def.}}{=} & (\text{fun list} \rightarrow \text{match list with} \\ & \quad [] \rightarrow [] \\ & \quad | (x,y)::\text{tail} \rightarrow (y,x)::(\text{flip tail})) (a,b)::l \\ &= \text{match } (a,b)::l \text{ with} \\ & \quad [] \rightarrow [] \\ & \quad | (x,y)::\text{tail} \rightarrow (y,x)::(\text{flip tail}) \\ &= (b,a)::(\text{flip } l) \end{aligned}$$

Daraus folgt:

$$\begin{aligned} &\text{flip}(\text{flip list}) \\ &= \text{flip}(\text{flip } ((x',y')::\text{list}')) \\ &= \text{flip } ((y',x')::(\text{flip list}')) \\ &= (x',y')::(\text{flip } (\text{flip list}')) \\ \stackrel{\text{I.V.}}{=} & (x',y')::\text{list}' \\ &= \text{list} \end{aligned}$$