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Virtual Machines

Summer Semester 2007

Exercise sheet 10

Deadline: 17 July 2007 12:00

Aufgabe 1: Unification: Occur check

Unify t and s and compare the costs for unification with and without occur check.

 $t \equiv p(X_1, \dots, X_n)$ $s \equiv p(f(X_0, X_0), f(X_1, X_1), \dots, f(X_{n-1}, X_{n-1}))$

Aufgabe 2: Code Generation for WiM

Consider the following program P: rev1([], R, R) :- . rev1([X|L], R, A) :- rev1(L, [X|R], A). reverse(L, R) :- rev1(L, [], R).

?- reverse(X, [4, 2, 1]).

- (a) Translate P to WiM code (with optimization) and point them out.
- (b) Execute the WiM code showing the sequence of (sub-)goals that are called and monitor how stack and heap develop after each of these goals has been processed.

Aufgabe 3: Prolog-arithmetics

In Prolog the arithmetic operators +, -, * and / are interpreted as common term constructors. In order to implement an efficient arithmetics, Prolog is extended by the type Integer and the predicate is/2. Thus, the Goal "X is t" arithmetically evaluates the term t and unifies the result of the evaluation with X.

- a) The WiM should be extended by the data type **Integer** and (analogously to atoms) by the abstract instructions for unification of Integers.
- b) Define an abstract instruction eval, which pops the term t, pointing to the topmost stack cell, from the stack, arithmetically evaluates the term and pushes the result on the stack. If t contains unbounded variables, atoms or structures not built with +, -, * or /, the instruction eval should fail and raise an error.
- c) Analogously to MaMa define abstract instructions for performing arithmetical operations on the stack. Specify a code generation function code_I to handle arithmetical expressions on the right-hand side of an *is*-Goal.
- d) Specify a code scheme for goals of the form "X is t".
- e) Generate code_PR for the following predicate:

len([], 6-2*3) := .len([X|R], L) := len(R, L1), L is L1 + 1. 10 Punkte

2 Punkte

18 Punkte