

# Compiler Construction & Virtual Machines

## Exercise Sheet 1

Deadline: 23. April 2008, 12:00, during the lecture or in room 02.07.053

### Exercise 1: Code generation

8 Points

Consider the following instruction sequence.

```
z = 1;
while (n > 0) {
    j = 1;
    y = x;
    while (2 * j <= n) {
        y = y * y;
        j = j * 2;
    }
    z = y * z;
    n = n - j;
}
```

- What does the instruction sequence compute?
- Generate CMA code for the instruction sequence.  
Use the address environment  $\rho = \{n \mapsto 0, j \mapsto 1, x \mapsto 2, y \mapsto 3, z \mapsto 4\}$  !

### Exercise 2: Registers

12 Points

We extend the CMA machine by adding an unbounded number of registers  $R_0, R_1, \dots$ . To improve efficiency, expressions are evaluated by storing intermediate values in registers instead of on the stack. For example, to evaluate  $x * y + 2$  and to store the result in  $R_1$ , we first put the address of  $x$  in  $R_1$ , put  $S[R_1]$  in  $R_1$ , put the address of  $y$  in  $R_2$ , put  $S[R_2]$  in  $R_2$ , put  $R_1 + R_2$  in  $R_1$ , put 2 in  $R_2$  and then put  $R_1 + R_2$  in  $R_1$ .

- Choose a set of new CMA instructions for doing this translation.
- Give the translation scheme for evaluation of expressions and assignment statements. For this purpose, extend the functions  $code_L$  and  $code_R$  to now take an additional argument  $R$  which is the register in which to store the result of evaluation. We assume the invariant that all registers  $R_j$  with  $j \geq i$  are free.